

Escape the Matrix: Graphical Reasoning and Minimal Axioms for Quantum Circuits

Colin Blake

PhD Day Edition 2025

May 6, 2025



Equational Theories via Music

Completeness of Equational Theories

Minimality of Axioms

Minimal Axioms for Quantum Fragments

Example Axioms: Qubit Clifford

Conclusion

Motivation: The Matrix Problem

What happens if you try to verify a *20-gate* quantum circuit by hand?

- You multiply 20 matrices, each $2^n \times 2^n$, by hand
- Errors multiply and it becomes intractable for even 5 qubits

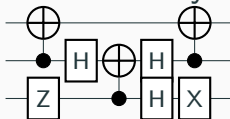
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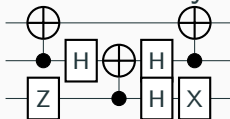
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- Rewrite rules on wires & nodes
- Local transformations, no full-blown matrix multiplication
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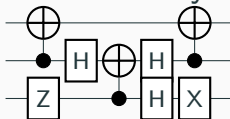
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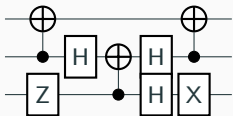


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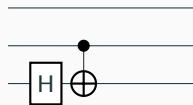
Circuit Optimization with Rewrites

Goal: Simplify a complex quantum circuit while preserving its function.

Before:



After:



Such rewrites come from a set of *equational rules* (axioms).

Equational Theory: Musical Analogy

- **Generators:** Musical note durations (whole, half, quarter, ...)
- **Equations:** Duration identities, e.g. $\text{half} + \text{half} = \text{whole}$
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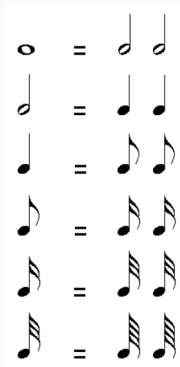
Note Duration Example

Original

half note + half note

Rewritten

whole note



A small, minimal rule delivers powerful rewriting capabilities.

What is Completeness?

A set of circuit equations is **complete** if:

- *Any* two circuits implementing the same unitary can be transformed into each other by these rules.
- Guarantees no missing identities: every true equality is derivable.

Why it matters: Ensures we can automate verification of equivalence.

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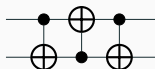
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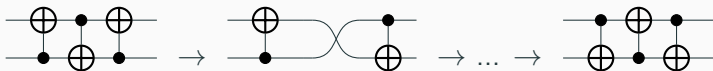
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Completeness Example: SWAP Circuits

Two implementations of a SWAP gate:



A complete theory provides a sequence of rewrites:



proving semantic equivalence without matrices.

What is Minimality?

- **Minimality:** No axiom is redundant.
- Removing any single rule breaks completeness: some equalities become unprovable.
- Think of each axiom as a "+1 leg" supporting the theory table.

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Proving Independence via Interpretations

- To show axiom X is necessary:
 - Construct an *alternate interpretation* where all other axioms hold,
 - But axiom X fails (gives different semantics).
- This counter-model demonstrates X cannot be derived from the rest.

Example: Interpretation which counts the parity of some generators.

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- **Qubit Clifford** : Core operations. Widely used for error correction and efficient classical simulation. It uses phases $\pm 1, \pm i$ (complex roots).
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Equational Theory for Qubit Clifford Circuits

$$\omega^{\otimes 8} = \boxed{} \quad (\omega^8) \quad \boxed{H} \boxed{H} = \text{---} \quad (H^2)$$

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$$\boxed{H} \boxed{S} \boxed{H} = \omega \boxed{S^\dagger} \boxed{H} \boxed{S^\dagger} \quad (E)$$

$$\begin{array}{c} \bullet \\ \oplus \end{array} \boxed{S} \begin{array}{c} \bullet \\ \oplus \end{array} = \boxed{S} \quad (Cs) \quad \begin{array}{c} \bullet \\ \oplus \end{array} \begin{array}{c} \oplus \\ \bullet \end{array} = \text{---} \quad (B)$$

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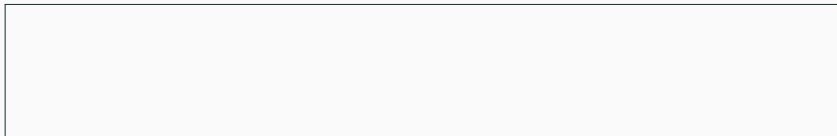
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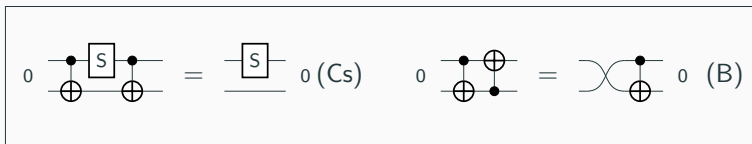
The diagram shows an equality between two quantum circuit fragments. On the left, a CNOT gate (represented by a dot on the top wire and a circle with a plus sign on the bottom wire) is followed by a SWAP gate (represented by a box labeled 'S'). On the right, a SWAP gate is followed by a CNOT gate. The two fragments are separated by an equals sign. The entire diagram is enclosed in a rectangular box.

$$0 \quad \begin{array}{c} \bullet \\ \oplus \end{array} \begin{array}{c} \bullet \\ \oplus \end{array} \begin{array}{c} \text{S} \\ \oplus \end{array} = \begin{array}{c} \text{S} \\ \oplus \end{array} \begin{array}{c} \bullet \\ \oplus \end{array} \begin{array}{c} \bullet \\ \oplus \end{array} 0(C_S)$$

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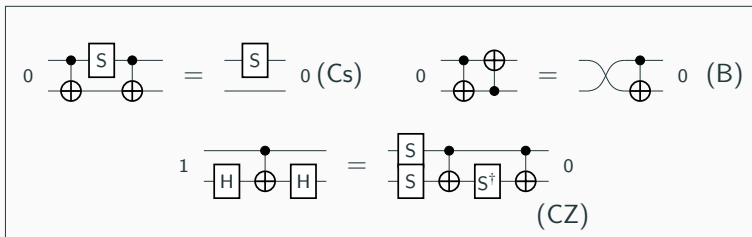
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Comparison of Axiom Set Sizes

Fragment	Previous Theory	Our Theory
Qubit Clifford	15	8 (minimal for n -qubits)
Real Clifford	16	10 (minimal for n -qubits)
Qutrit Clifford	17	11 (minimal for 2-qutrits, conjectured for n)
Clifford+T	18	11 (minimal for 1-qubit, simplified for 2, no result for n)

Summary

- Graphical equational theories let us *reason* and *optimize* circuits visually.
- Our rule sets are *complete* (all true equivalences provable) and are *minimal* in the sense that we can't derive some equations from the others.
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